Compilers

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Lecture#3

Implementation of Lexical Analysis

Outline

- Specifying lexical structure using regular expressions
- Finite automata
  - Deterministic Finite Automata (DFAs)
  - Non-deterministic Finite Automata (NFAs)
- Implementation of regular expressions

  RegExp => NFA => DFA => Tables

Notation

- How regular expressions are used to construct a full lexical specification on a programming language?

- Ay least one: A^+ ≡ AA^*
- Union: $A \mid B$ ≡ $A + B$
- Option: $A + \varepsilon$ ≡ $A?$
- Range: ‘a’+’b’+…+’z’ ≡ [a-z]
- Excluded range:
  complement of [a-z] ≡ [^a-z]

Regular Expressions in Lexical Specification

- A specification for the predicate

  $s \in L(R)$?????

- What we want to do?
- We want to know whether a given string is in this language or not.
- But a yes/no answer is not enough!
- Instead: partition the input string into tokens. Each one of these tokens is in $L(R)$.

- We adapt regular expressions to this goal

Regular Expressions => Lexical Spec. (1)

Steps to construct a full lexical specification from regular expressions

1. Write a rexp for the lexemes of each token
   - Number = digit +
   - Keyword = 'if' + 'else' + …
   - Identifier = letter (letter + digit)*
   - OpenPar = '('
   - …

Regular Expressions => Lexical Spec. (2)

2. Construct R, matching all lexemes for all tokens

R = Keyword + Identifier + Number + …

= R_1 + R_2 + …

- The union of all the regular expressions forms the lexical specification of a language?

Regular Expressions => Lexical Spec. (3)

3. Let input be $x_1 \ldots x_n$
   For $1 \leq i \leq n$ check
   $$x_1 \ldots x_i \in L(R)$$

4. If success, then we know that
   $$x_1 \ldots x_i \in L(R_j) \text{ for some } j$$

5. Remove $x_1 \ldots x_i$ from input and go to (3)

Repeat until the input string is empty (the entire input is analyzed)

Ambiguities (1)

- There are ambiguities in the algorithm
- How much input is used? What if
  - $x_1 \ldots x_i \in L(R)$ and also
  - $x_1 \ldots x_K \in L(R)$
- Rule: Pick longest possible string in $L(R)$
  - The "maximal munch"

Ambiguities (2)

- Which token is used?

What if

- $x_1 \ldots x_i \in L(R_j)$

and also

- $x_1 \ldots x_i \in L(R_k)$

- Rule: use rule listed first (j if $j < k$)
  - Treats "if" as a keyword, not an identifier

Error Handling

- What if
  No rule matches a prefix of input ?
- Problem: Can't just get stuck …Compilers should do good error handling.
- Solution:
  - Write a rule matching all "bad" strings

ERROR=[all strings not in the lexical specification]

  - Put it last (lowest priority)

Summary

- Regular expressions provide a concise notation for string patterns
- Use in lexical analysis requires small extensions
  - To resolve ambiguities
  - To handle errors
- Good algorithms known
  - Require only single pass over the input
  - Few operations per character (table lookup)

Finite Automata

- Regular expressions = specification
- Finite automata = implementation mechanism for regular expressions

- A finite automaton consists of
  - An input alphabet Σ
  - A finite set of states S
  - A start state n
  - A set of accepting states $F \subseteq S$
  - A set of transitions $\text{state} \rightarrow^{\text{input}} \text{state}$

Finite Automata

- Transition

$$s_1 \rightarrow^a s_2$$

- Is read

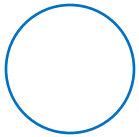
  In state s_1 on input "a" go to state s_2

- If end of input and in accepting state => accept

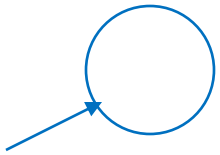
- Otherwise => reject

Finite Automata State Graphs

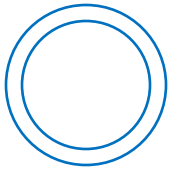
A state



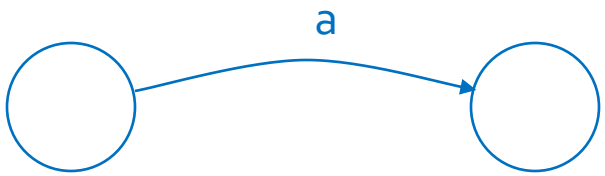
The start state



An accepting state

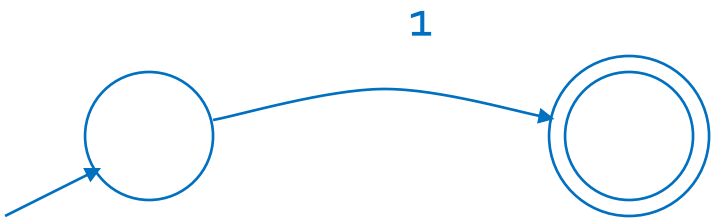


A transition



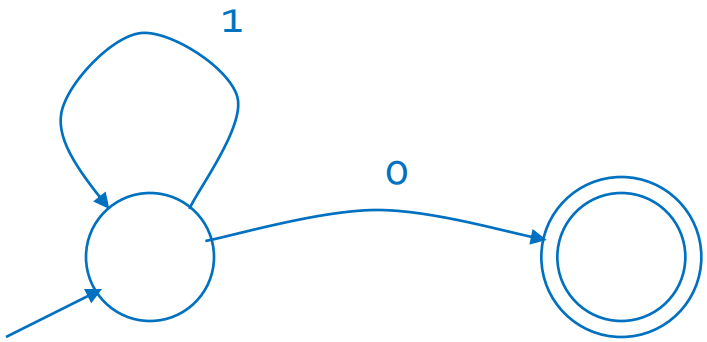
A Simple Example

- A finite automata that accepts only "1"



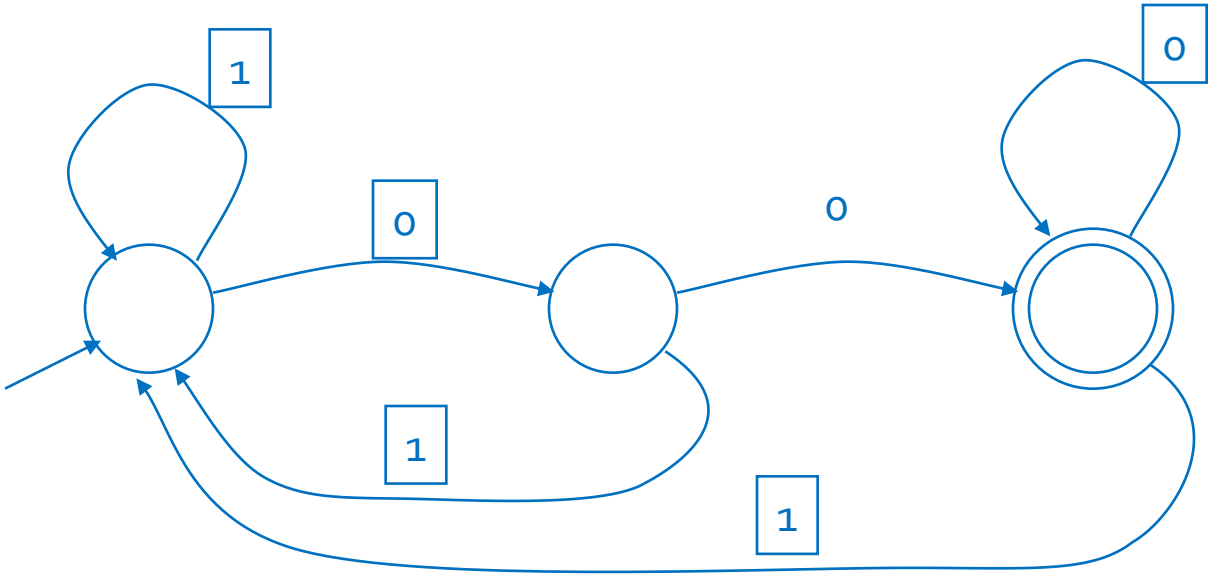
Another Simple Example

- A finite automaton accepting any number of 1's followed by a single 0
- Alphabet: {0,1}



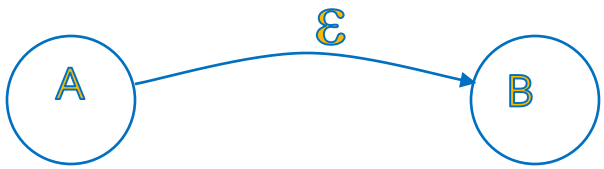
Another Example

- Alphabet {0,1}
- What language does this recognize?



Epsilon Moves

- Another kind of transition: ε-moves



- Machine can move from state A to state B without reading input
- Think of ε-moves as a kind of free move.

Deterministic and Non-deterministic Automata

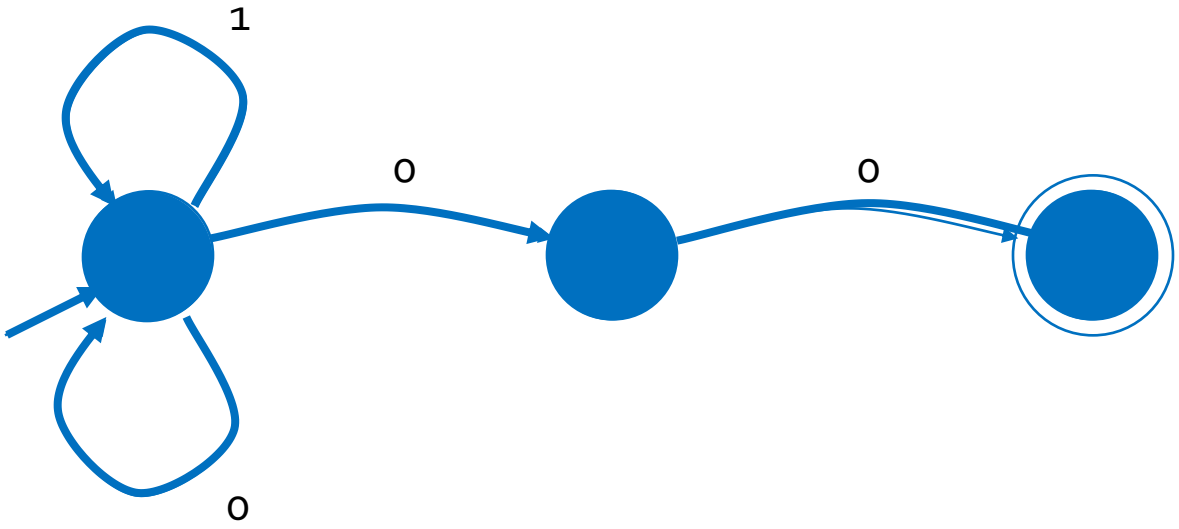
- Deterministic Finite Automata (DFA)
  - One transition per input per state
  - No ε-moves
- Nondeterministic Finite Automata (NFA)
  - Can have multiple transitions for one input in a given state
  - Can have ε-moves

Execution of Finite Automata

- A DFA can take only one path through the state graph
  - The path is completely determined by the input (The machine has no choice)
- NFAs can choose
  - Whether to make ε-moves
  - Which of multiple transitions for a single input to take

Acceptance of NFAs

- An NFA can get into multiple states



- Input: 1 0 0

Rule: NFA accepts if it can get to a final state

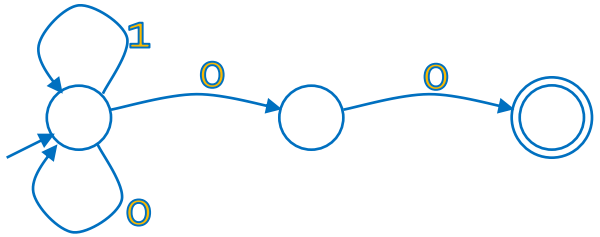
NFA vs. DFA (1)

- NFAs and DFAs recognize the same set of languages (regular languages)
- DFAs are faster to execute
  - There are no choices to consider

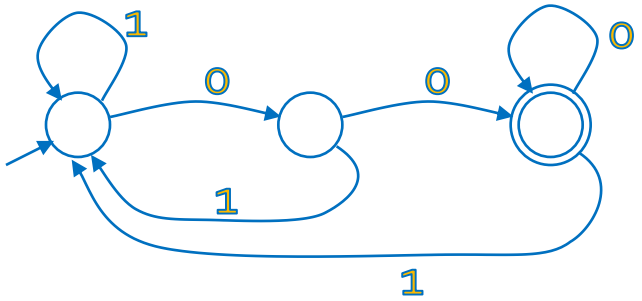
NFA vs. DFA (2)

- For a given language NFA can be simpler than DFA

NFA



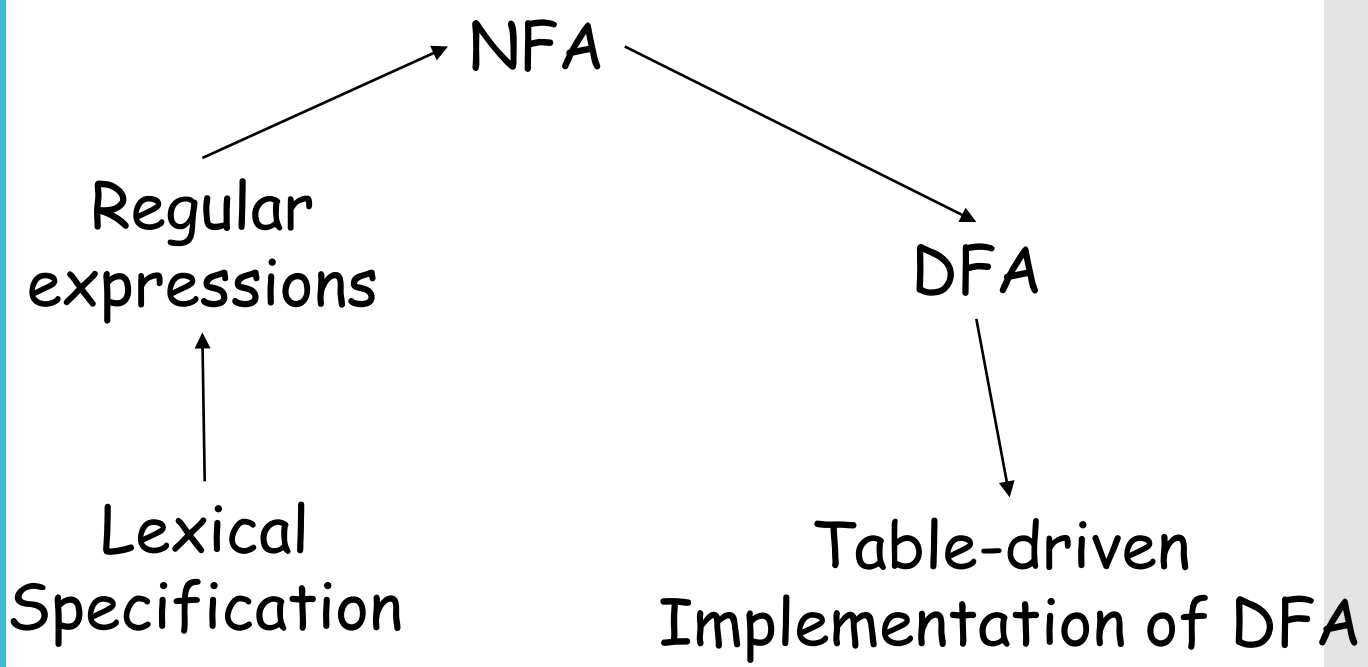
DFA



DFA can be exponentially larger than NFA

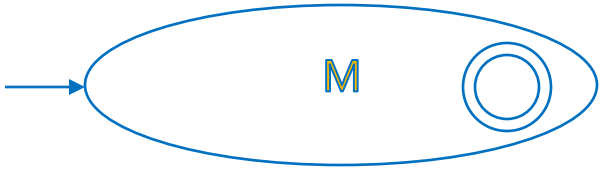
Regular Expressions to Finite Automata

- High-level sketch



Regular Expressions to NFA (1)

- For each kind of rexp, define an NFA
  - Notation: NFA for rexp M

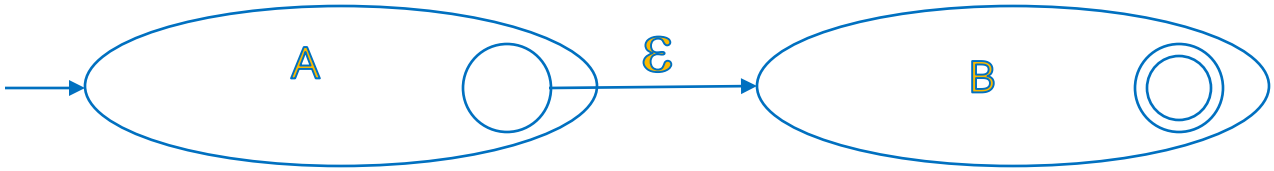


- For ε

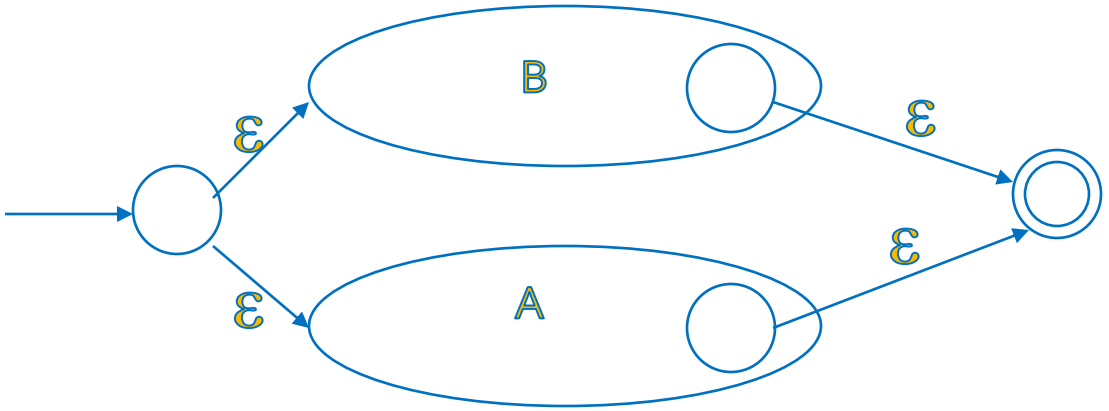
- For input a

Regular Expressions to NFA (2)

- For AB

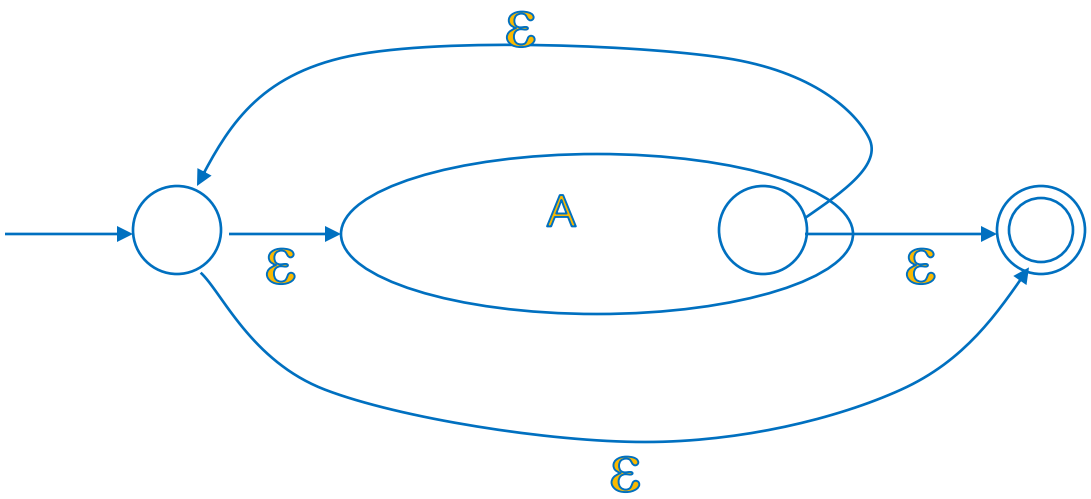


- For *A + B*



Regular Expressions to NFA (3)

- For A*

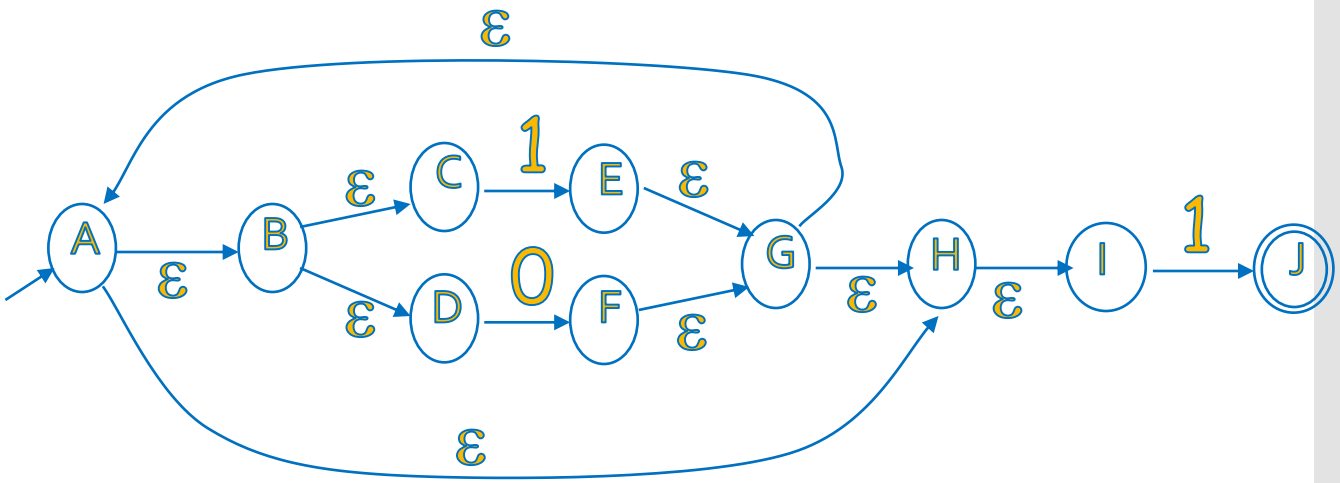


Example of RegExp -> NFA conversion

- Consider the regular expression

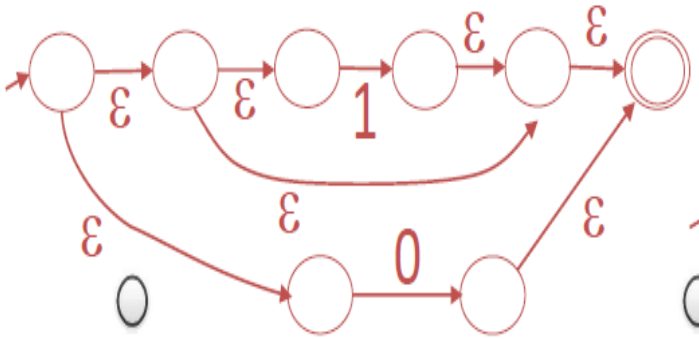
  (1+0)*1

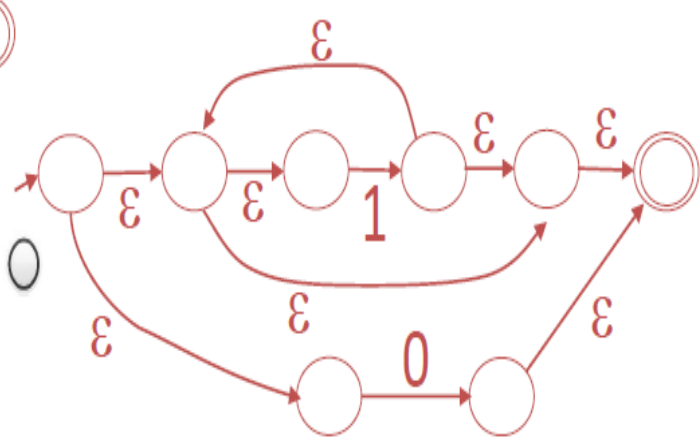
- The NFA is

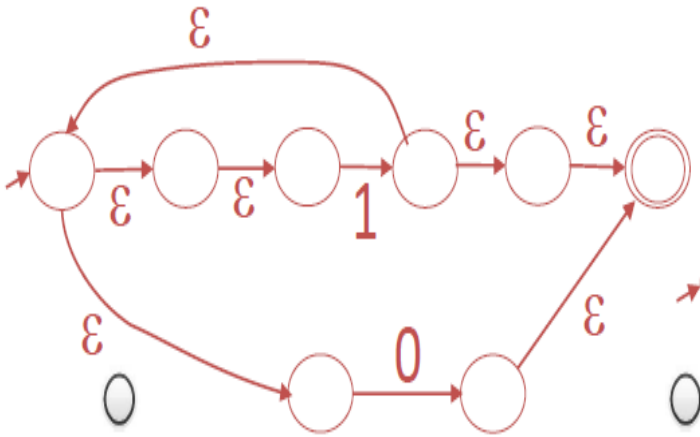


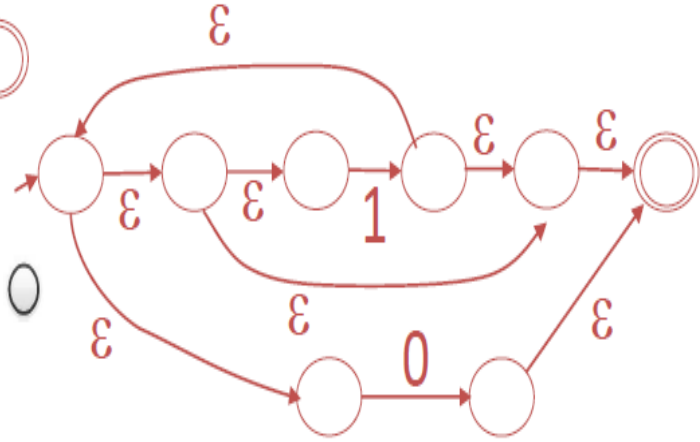
Quiz

Choose the NFA that accepts the following regular expression: 1* + 0

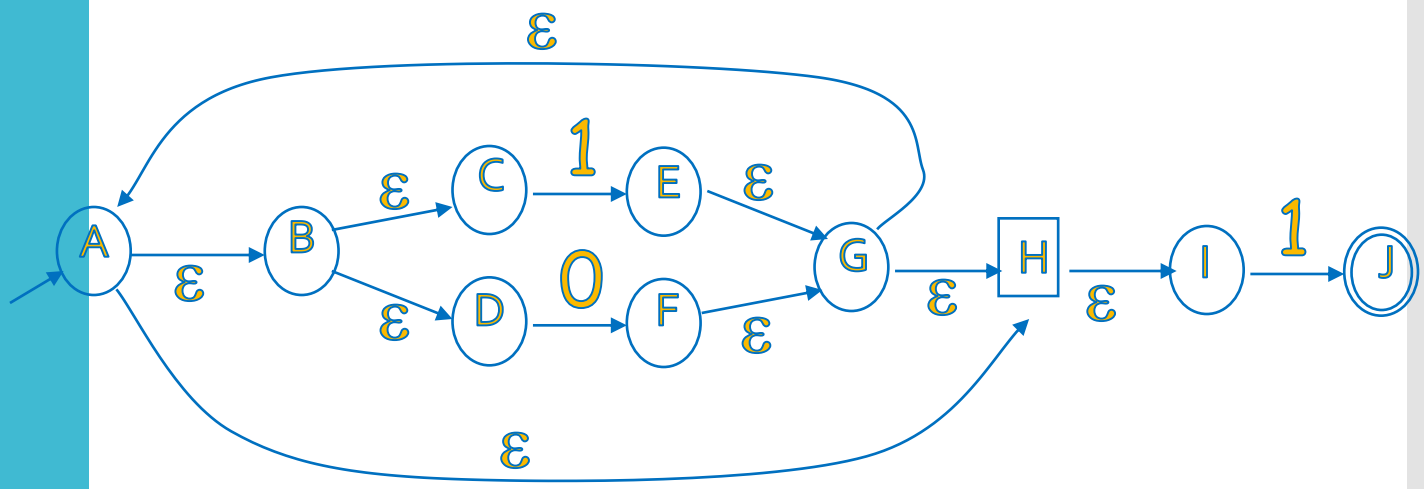








NFA to DFA.



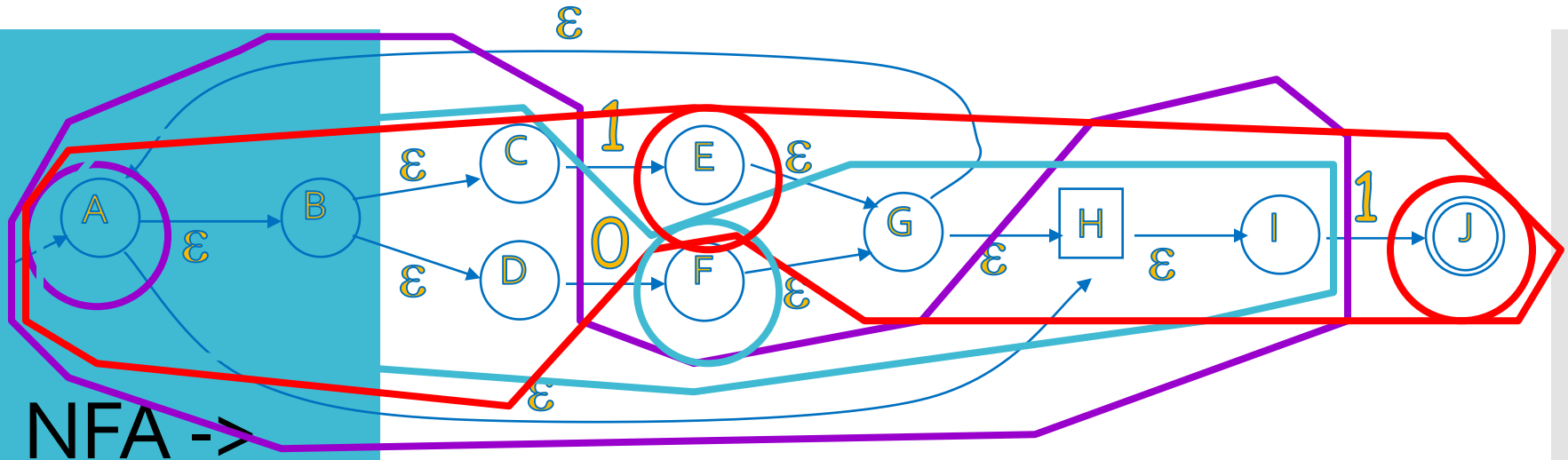
NFA to DFA: *The Trick*

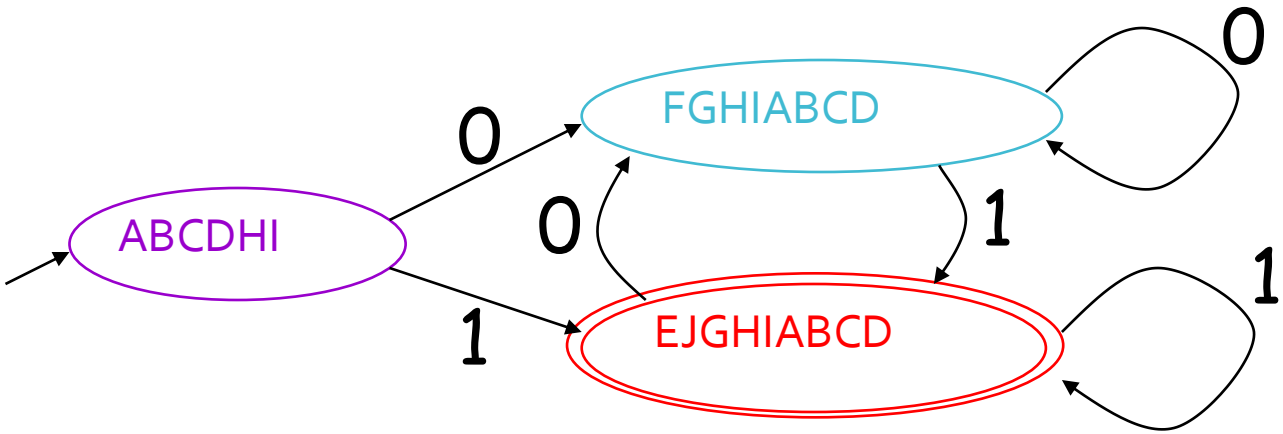
- Simulate the NFA
- Each state of DFA
  = a non-empty subset of states of the NFA
- Start state
  = the set of NFA states reachable through ε-moves from NFA start state
- Add a transition $S \rightarrow^{a} S'$ to DFA iff
  - S' is the set of NFA states reachable from any state in S after seeing the input a, considering ε-moves as well

NFA to DFA. Remark

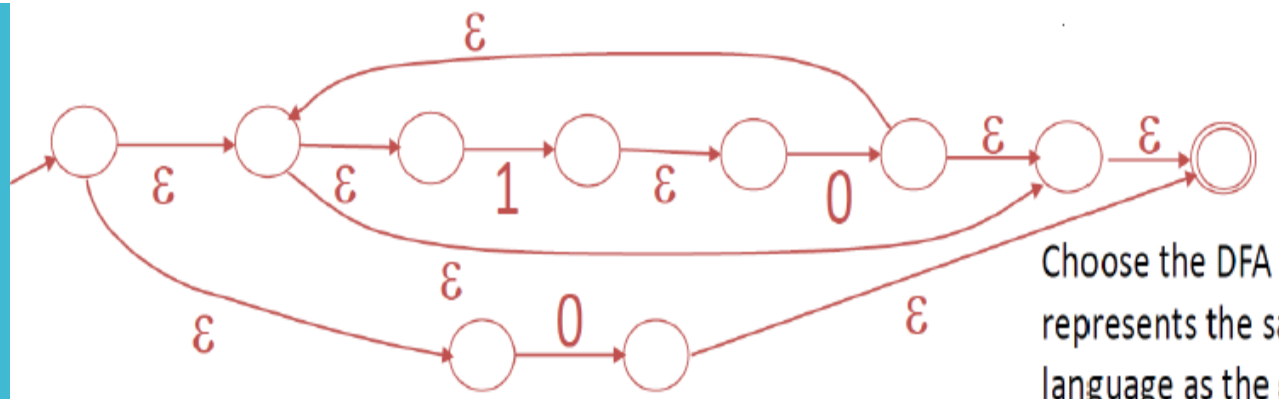
- An NFA may be in many states at any time
- How many different states ?
- If there are N states, the NFA must be in some subset of those N states
- How many subsets are there?
  - $2^N - 1$ = finitely many

NFA -> DFA Example

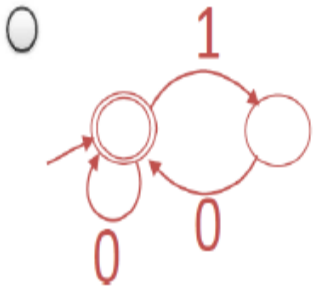


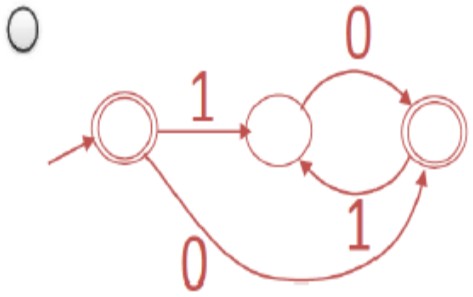


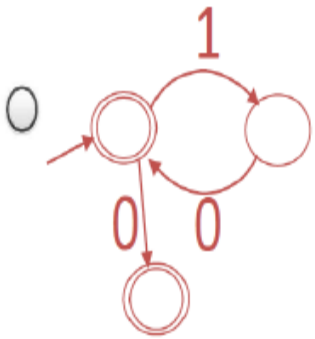
QUIZ

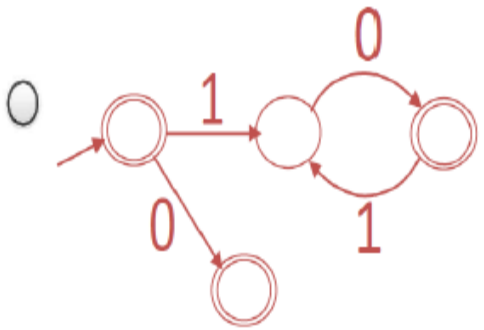


Choose the DFA that represents the same language as the given NFA





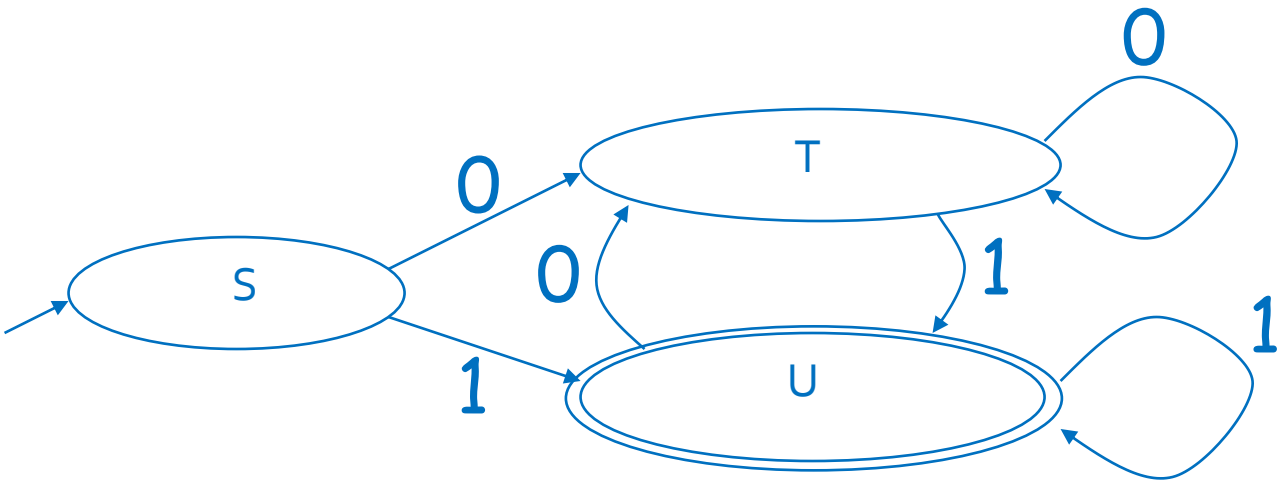




Implementation

- A DFA can be implemented by a 2D table T
  - One dimension is "states"
  - Other dimension is "input symbol"
  - For every transition $S_i \rightarrow^a S_k$ define $T[i,a] = k$
- DFA "execution"
  - If in state S_i and input a, read $T[i,a] = k$ and skip to state S_k
  - Very efficient

Table Implementation of a DFA



	0	1
S	T	U
T	T	U
U	T	U

Implementation (Cont.)

- NFA -> DFA conversion is at the heart of tools such as flex
- But, DFAs can be huge
- In practice, flex-like tools trade off speed for space in the choice of NFA and DFA representations
- DFA : faster, less compact
- NFA: slower, consumes less memory

Thanks